The Multisecretary Problem with many types

Omar Besbes, Yash Kanoria, Akshit Kumar

Columbia Business School

Problem Introduction ••••					
» Multi-se	ecretary Pr	roblem			

Problem Statement

Given a sequence of *T* secretaries and a hiring budget *B*, a decision maker (DM) wants to hire the top *B* secretaries in terms of their ability.

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» Multi-secretary Problem

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Details and Assumptions

* The secretaries arrive in an online fashion.

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Details and Assumptions

- * The secretaries arrive in an online fashion.
- * The DM makes **irrevocable** *hire* or *reject* decisions.
- * The abilities (types) of the secretaries are drawn **independently** from a **common** and **known** distribution *F* over [0, 1].

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» Gotta catch'em all some







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DM	t = 1	$\frac{1}{t=2}$	t = 3	t=4	$\frac{1}{t=5}$	t = 6]	



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DM	t = 1	$\frac{1}{t=2}$	t = 3	$\frac{1}{t=4}$	$\frac{1}{t=5}$	$\frac{1}{t=6}$]	



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DM	t = 1	$\frac{1}{t=2}$	t=3	$\frac{1}{t=4}$	$\frac{1}{t=5}$	t = 6	t = 7



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Online Policy

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Failure of CE

Conclusion

- » Motivation and Related Applications
 - * Online Knapsack Problem [Kleywegt and Papastavrou, 1998]
 - * Unit weights and *i.i.d* rewards

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 - * Single-leg Revenue Management [Talluri et al., 2004]
 - * Selling horizon *T*, flight capacity *B* and types are different fare classes.

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 - * Order Fulfillment Problem [Lei et al., 2018]

» Motivation and Related Applications

- * Online Knapsack Problem [Kleywegt and Papastavrou, 1998]
 - * Unit weights and *i.i.d* rewards
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 - Selling horizon T, flight capacity B and types are different fare classes.
- * Order Fulfillment Problem [Lei et al., 2018]

Spatially distributed demand 2 WHs with total inventory *T* Unit demand at each time Fulfill from one of the WHs Inventory depletes by 1 min total matching cost



Reduces to multi-secretary problem



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Continuous Distributions "Many Types" \uparrow \uparrow \uparrow \downarrow \downarrow Regret = $\Theta(\log T)$ [Bray, 2022], [Lueker, 1998]





Continuous Distributions "Many Types" Regret = $\Theta(\log T)$ [Bray, 2022], [Lueker, 1998] Need density to be bounded below



















» Punchline and Overview



Rarity of types / Shape of the density (β)



» Punchline and Overview



Drivers of Regret

Rarity of types / Shape of the density (β)





Rarity of types / Shape of the density (β)



» Punchline and Overview



Rarity of types / Shape of the density (β)

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» Punchline and Overview



- Shape is a fundamental driver of regret
- Dealing with gaps is an algorithmic challenge.
- New principle: Conservatism wrt gaps
- * If close to the gap, use the gap.

Rarity of types / Shape of the density (β)

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» First step towards unified approach









- * Consider a continuous distribution with finite many points at which the density is zero.
- * β characterizes the rate of mass accumulation around the points of zero density.
- * Consider *m* such that f(m) = 0, then *F* is $(\beta, 1)$ -clustered if

$$|F(m \pm \delta) - F(m)| \ge \delta^{\beta+1}$$

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		Continuous Distributions		
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» $(m{eta}, 1)$ -clustered distributions

Universal Lower Bound

For every $eta\in[0,\infty)$, there exists a distribution F_eta such that

$$\sup_{B \in [T]} \mathbb{E}_{F_{\beta}} \left[\mathsf{Regret} \right] = \begin{cases} \Omega \left(\log T \right), & \beta = 0, \\ \Omega \left(T^{\frac{1}{2} - \frac{1}{2(1+\beta)}} \right), & \beta > 0. \end{cases}$$

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	Continuous Distributions		

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Regret of Certainty Equivalent

If *F* is a $(\beta, 1)$ -clustered distribution, then for all $B \in [T]$,

$$\mathbb{E}\left[\mathsf{Regret}(\mathsf{CE})\right] = \begin{cases} \mathcal{O}\left(\log T\right), & \beta = 0, \\ \mathcal{O}\left(T^{\frac{1}{2} - \frac{1}{2(1+\beta)}}\right), & \beta > 0. \end{cases}$$

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» Certainty Equivalent Control

Continuous Distributions

For $(m{eta}, 1)$ -continuous distributions



- * Let B_t be the remaining budget at time t
- * Compute the budget ratio $br_t = \frac{\text{Remaining Budget}}{\text{Remaining Time}} = \frac{B_t}{T-t}$
- * Define a quantile threshold $p_t^{ce} = 1 br_t$
- * Define a ability threshold $\gamma_t^{ce} = \mathit{F}^{-1}(\mathit{p}_t^{ce})$
- * hire \iff $heta_t \geq \gamma_t^{ce}$

			Continuous Distributions			
$\gg (eta,1)$	-clustered	l distributions			Brief Sum	Imar

- * Relax the assumption of density being positively lower bounded for continuous distribution.
- * The lower bound along with the upper bound show that β (shape of the density) is a fundamental driver of regret.
- * Entire spectrum of regret scaling is possible: $\Theta(\log T)$ or $\Theta(T^{\alpha})$ where $\alpha \in (0, \frac{1}{2})$.
- * A simple heuristic like CE is optimal.

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» Gaps in the distribution





» Gaps in the distribution





» Gaps in the distribution





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				Gaps in distribution		
» $(eta, \epsilon$	$(\mathbf{c_0})$ -cluster	ed distributio	J		Exar	nples



 $f(\theta)_{\uparrow}$

				Gaps in distribution ○●○		
» ($eta,arepsilon$	o)-cluster	ed distributio	٦		Exar	nples





 $\text{Gap} \equiv \text{intervals}$ of positive length with zero mass

mass cluster \equiv interval with positive mass



 $f(\theta)_{\uparrow}$





 $\text{Gap} \equiv \text{intervals}$ of positive length with zero mass

mass cluster \equiv interval with positive mass

 $\beta = 0$ (mass accumulation around gaps)



 $f(\theta)_{\uparrow}$



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Examples

 $\text{Gap} \equiv \text{intervals}$ of positive length with zero mass

mass cluster \equiv interval with positive mass

 $\beta = 0$ (mass accumulation around gaps)







 $\text{Gap} \equiv \text{intervals}$ of positive length with zero mass

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 $|F(m + \delta) - F(m)| \ge \delta$ on the same mass cluster





Examples

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mass cluster \equiv interval with positive mass

 $\beta = 0$ (mass accumulation around gaps)

 $|F(m + \delta) - F(m)| \ge \delta$ on the same mass cluster

 $\mu(\text{mass clusters}) \geq \varepsilon_0$





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 $Gap \equiv intervals$ of positive length with zero mass

mass cluster \equiv interval with positive mass

 $\beta = 0$ (mass accumulation around gaps)

 $|F(m + \delta) - F(m)| \ge \delta$ on the same mass cluster

 μ (mass clusters) $\geq \varepsilon_0$

For discrete distrbutions, $\beta = 0$, $\varepsilon_0 = \min_j p_j$

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mass clusters 1

gap

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 $f(\theta)_{\uparrow}$

 $Gap \equiv intervals$ of positive length with zero mass

mass cluster \equiv interval with positive mass

 $\beta = 1$ (mass accumulation around gaps)

 $|F(m + \delta) - F(m)| \ge \delta^2$ on the same mass cluster

 $\mu(\text{mass clusters}) \geq \varepsilon_0$

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» Certainty Equivalent Control

Continuous Distributions

n distribution Failure of LE CwG

For Bi-modal Uniform Distribution

Let B_t be the remaining budget at time t

$$\mathsf{Budget}\ \mathsf{Ratio} = rac{\mathsf{Remaining}\ \mathsf{Budget}}{\mathsf{Remaining}\ \mathsf{Time}} = rac{B_t}{T-t}$$

CE Quantile Threshold
$$= 1 - \frac{B_t}{T-t} \triangleq p_t^{ce}$$

Decision: hire $\iff \theta_t \ge F^{-1}(p_t^{ce})$



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Gaps in distribution

Failure of CE ○● Conclusio

» Failure of Certainty Equivalent Control

Regret Lower Bound

Insufficiency of Certainty Equivalent Control

Assume that $F = \text{Unif}([0, \frac{1}{4}] \cup [\frac{3}{4}, 1])$, for B = T/2, we have

 $\mathbb{E}\left[\mathsf{Regret}(\mathsf{CE})\right] = \Omega\left(\sqrt{\mathsf{T}}\right)$

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Remark

* Same scaling is achievable under a static threshold policy.

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Algorithmic Idea



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Algorithmic Idea



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Algorithmic Idea





Algorithmic Idea



If far from a gap, use the CE threshold

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Algorithmic Idea



If far from a gap, use the CE threshold



Algorithmic Idea



If far from a gap, use the CE threshold



Algorithmic Idea



If far from a gap, use the CE threshold If close to gap, use the gap as threshold



Algorithmic Idea



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Regret of CwG Policy

If F is a (β, ε_0) -clustered distribution, then

$$\mathbb{E}\left[\mathsf{Regret}(\mathsf{CwG})\right] = \begin{cases} \mathcal{O}\left((\log T)^2\right), & \beta = 0, \\ \mathcal{O}\left(\mathsf{poly}(\log T)T^{\frac{1}{2} - \frac{1}{2(1+\beta)}}\right), & \beta > 0 \end{cases}$$

If *F* is a discrete distribution, $\mathbb{E}[\text{Regret}(CwG)] = O(1/\varepsilon_0)$

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If *F* is a discrete distribution, $\mathbb{E} [\text{Regret}(CwG)] = O(1/\varepsilon_0)$

Remark

- * $F = \text{Unif}([0, \frac{1}{4}] \cup [\frac{3}{4}, 1])$, CwG ($\mathcal{O}((\log T)^2)$) outperforms CE ($\Omega(\sqrt{T})$).
- $\ast\,$ Matches the universal lower bound upto polylog factors $\,\Rightarrow\,$ CwG is near-optimal.

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				Conclusion • • • •
» Conclusio	ОП			

* *Motivation:* Discrete dist. and continuous dist. with bounded density are well studied but require different algorithmic ideas.

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» Conclusio	nc			

- * *Motivation:* Discrete dist. and continuous dist. with bounded density are well studied but require different algorithmic ideas.
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- * Challenges: Unified approach requires dealing with two potential drivers:
 - * Rarity of types around points of zero density (β): fundamental driver
 - * Presence of gaps: avoidable losses but need smarter algorithms



- * *Motivation:* Discrete dist. and continuous dist. with bounded density are well studied but require different algorithmic ideas.
- * *Research Ques:* How to unify and interpolate between them?
- * Challenges: Unified approach requires dealing with two potential drivers:
 - * Rarity of types around points of zero density (β): fundamental driver
 - * Presence of gaps: avoidable losses but need smarter algorithms
- * Algorithmic Innovation: Conservatism with respect to gaps can deal with a broader class of distributions.
 - * Recover the guarantees for discrete and bounded continuous distributions.

				Conclusion
» Open Pro	oblems			

* The multi-secretary problem is NRM for d = 1, how can we generalize these ideas to higher dimensions?

				Conclusion
» Open Pr	oblems			

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- * How to think of gaps in higher dimensions and how to deal with them?

				Conclusion OOO
» Open Pr	oblems			

- * The multi-secretary problem is NRM for d = 1, how can we generalize these ideas to higher dimensions?
- * How to think of gaps in higher dimensions and how to deal with them?
- * Extensions of NRM problems to instances with gaps.

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- * The multi-secretary problem is NRM for d = 1, how can we generalize these ideas to higher dimensions?
- * How to think of gaps in higher dimensions and how to deal with them?
- * Extensions of NRM problems to instances with gaps.
 - * Application: Order fulfillment problem with more than 2 warehouses.
 - * Gaps in spatial demand distributions in practical settings.

Thank you!



» References

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