Dynamic Resource Allocation

Algorithmic Design Principles and Spectrum of Achievable Performances

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Dynamic Resource Allocation is ubiquitous

Network Revenue Management

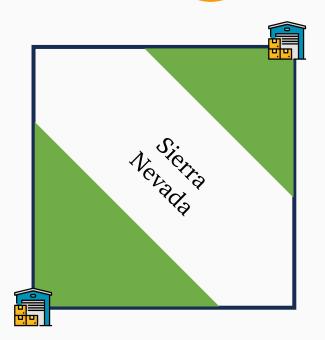


Budget Management in Ad auctions

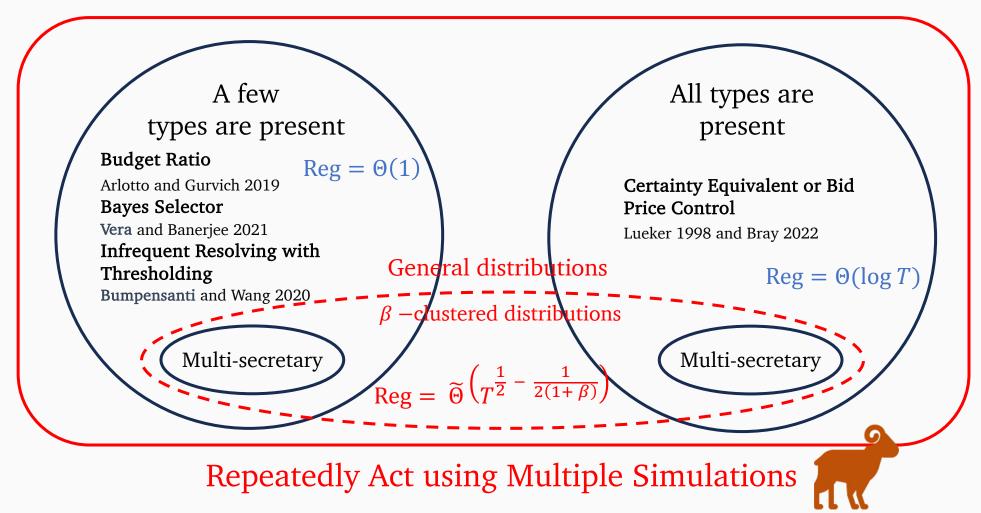
Bing ads AdWords

Order Fulfillment





In A Nutshell



Regret is the additive gap between the value of hindsight optimal problem and value of problem under some algorithm

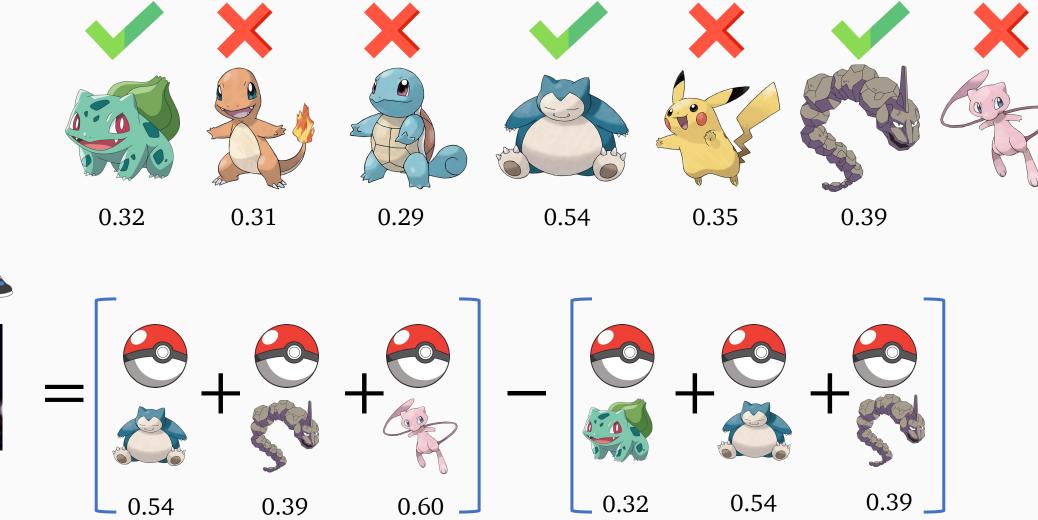
Multi-secretary problem

- Given a sequence of *T* secretaries and hiring budget *B*, a decision maker wants to hire the top *B* secretaries in terms of their ability.
- The secretaries arrive in an online fashion.
- The decision maker makes irrevocable hire or reject decisions.
- Assumption: The abilities (types) of the secretaries are drawn independently from a common and known distribution *F*.

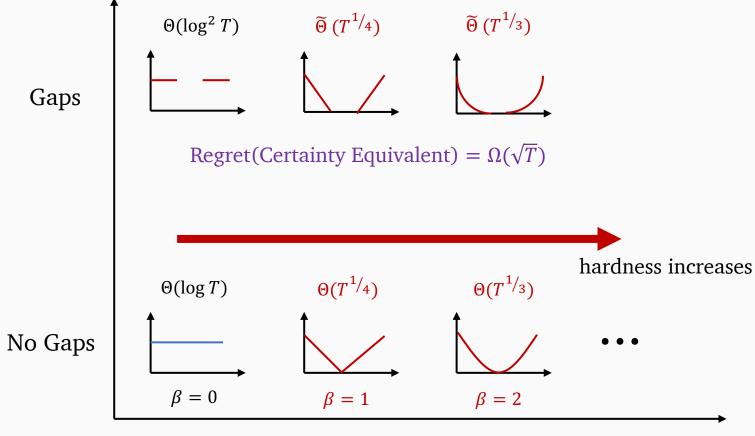
Multi-Pokémon Problem



Regret



Spectrum of Achievable Performances



Rarity of types (Shape of the distribution)

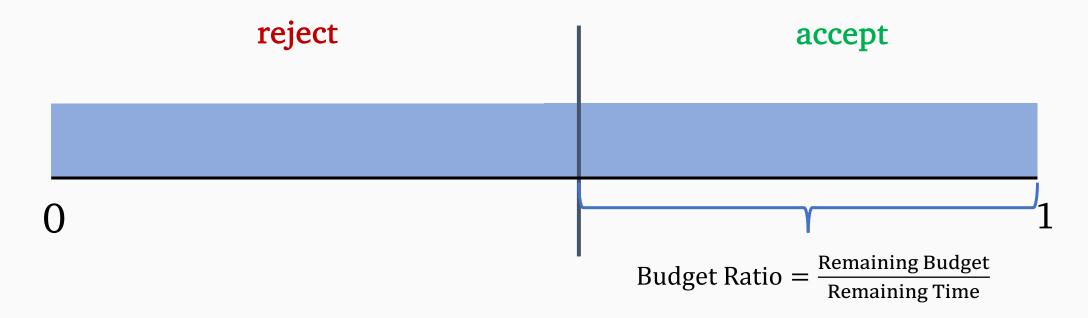
Distribution shape is fundamental driver of performance

Dealing with gaps is an algorithmic challenge

Conservativeness with respect to Gaps (CwG) principle enables near optimal performance

Use RAMS (Repeatedly Act using Multiple Scenarios) to operationalize CwG

Certainty Equivalent

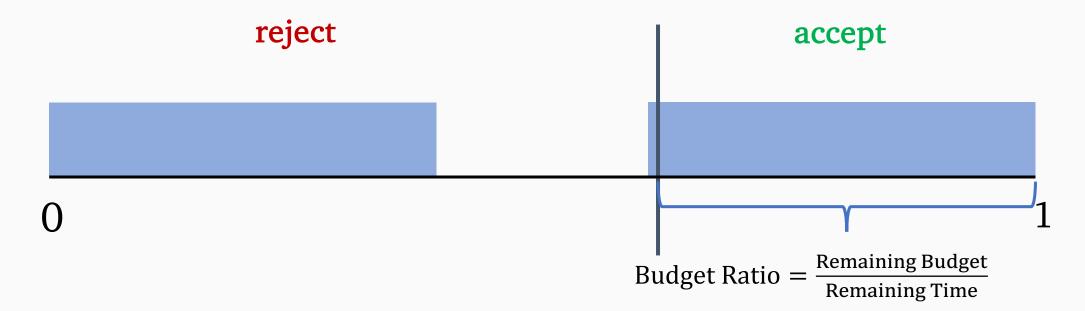


Uniform distribution (when all types are present)

$\operatorname{Regret}(\operatorname{CE}) = \Theta(\log T)$

(optimal regret scaling) Lueker 1998 and Bray 2019





Uniform distribution (when all types are present)

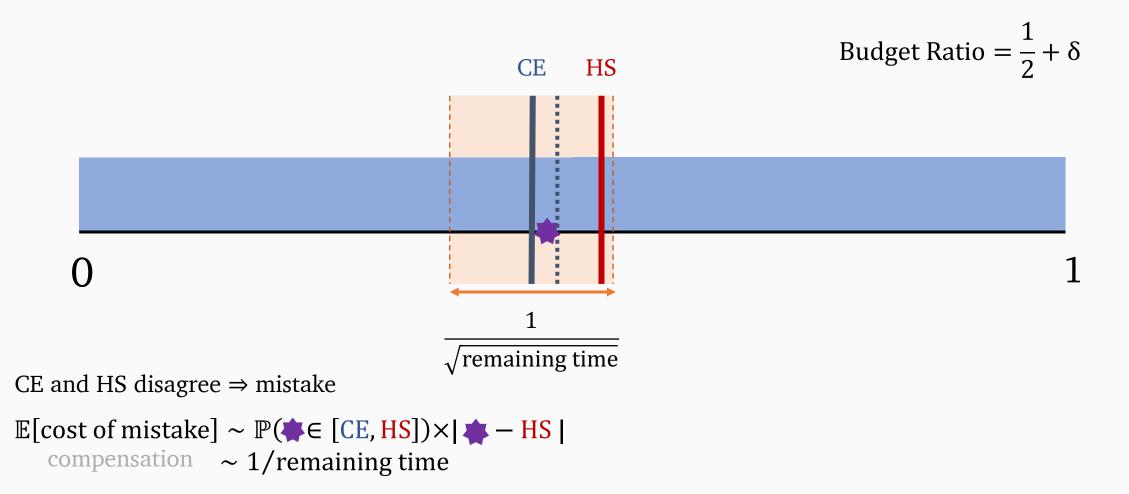
$\operatorname{Regret}(\operatorname{CE}) = \Theta(\log T)$

(optimal regret scaling) Lueker 1998 and Bray 2019 Bi-modal Uniform distribution (some types are absent)

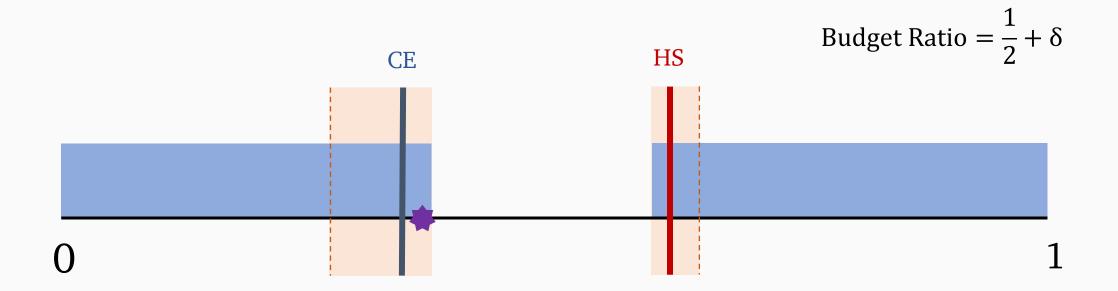
 $\operatorname{Regret}(\operatorname{CE}) = \Omega(\sqrt{T})$

(highly sub-optimal regret scaling)

Failure of Certainty Equivalent



Failure of Certainty Equivalent

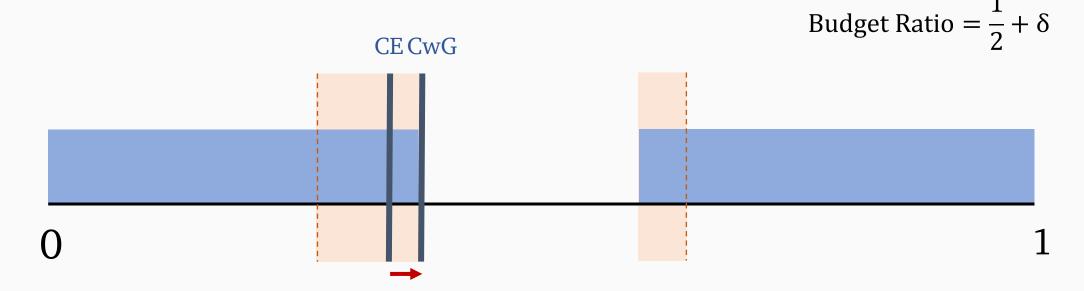


CE and HS disagree \Rightarrow mistake $\mathbb{E}[\text{cost of mistake}] \sim \mathbb{P}(\clubsuit \in [\text{CE}, \text{HS}]) \times | \clubsuit - \text{HS} |$

compensation $\sim 1/remaining time$

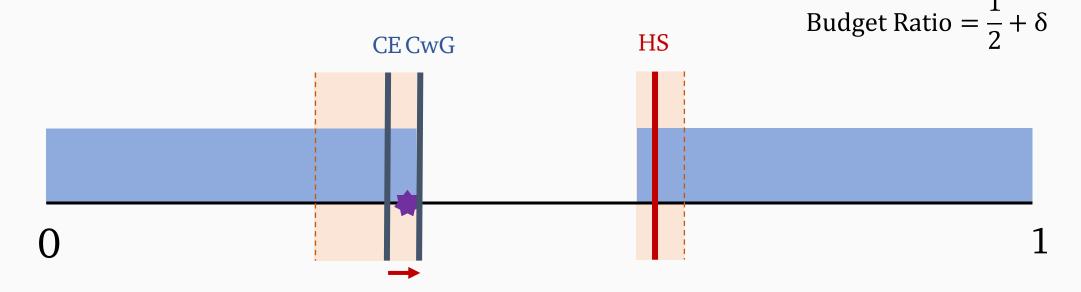
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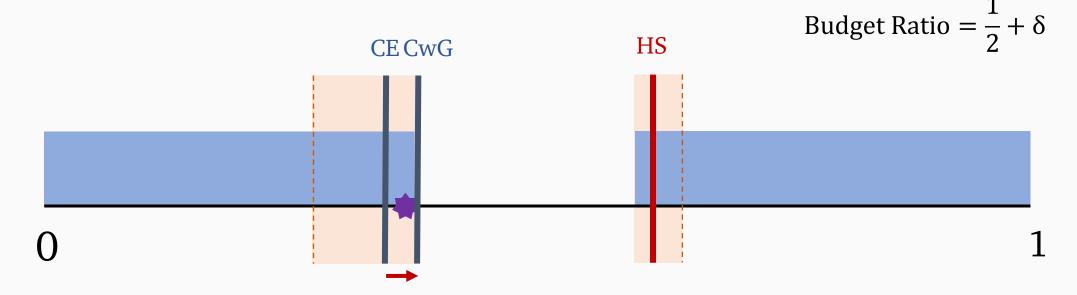
 $\mathbb{E}[\text{cost of mistake}] \sim \mathbb{P}(\clubsuit \in [\text{CE}, \text{HS}]) \times | \clubsuit - \text{HS} |$ compensation ~ $1/\sqrt{\text{remaining time}}$



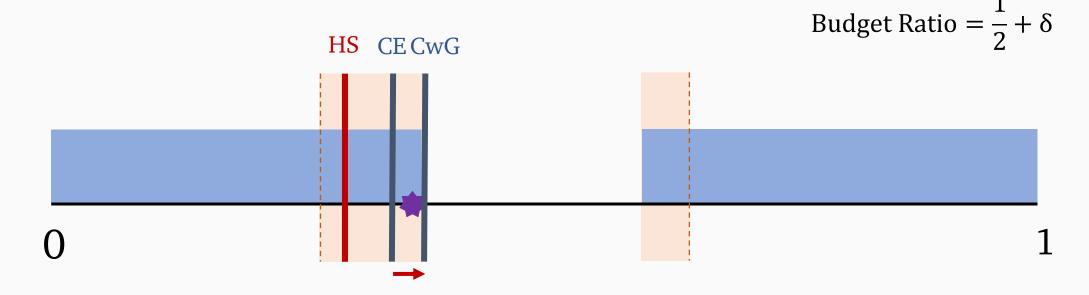
Conservativeness with respect to gaps principle

If the CE threshold is close to the gap, use the gap as the threshold. Otherwise use the CE threshold.



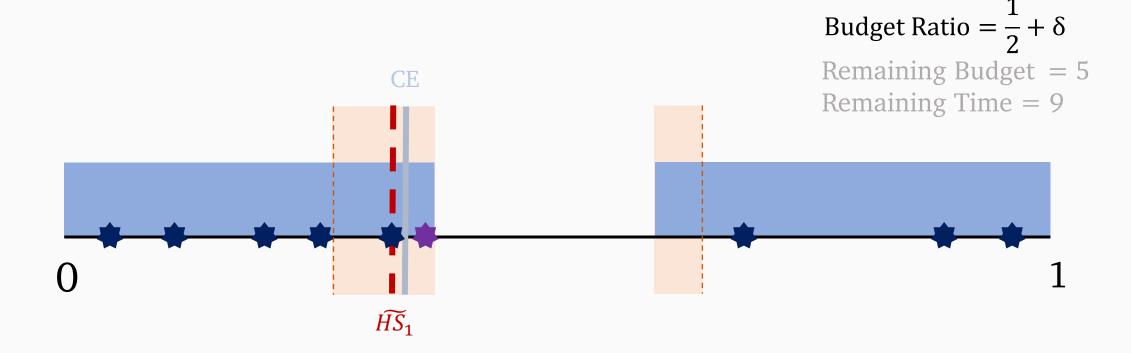


CwG and HS agree \Rightarrow no compensation required

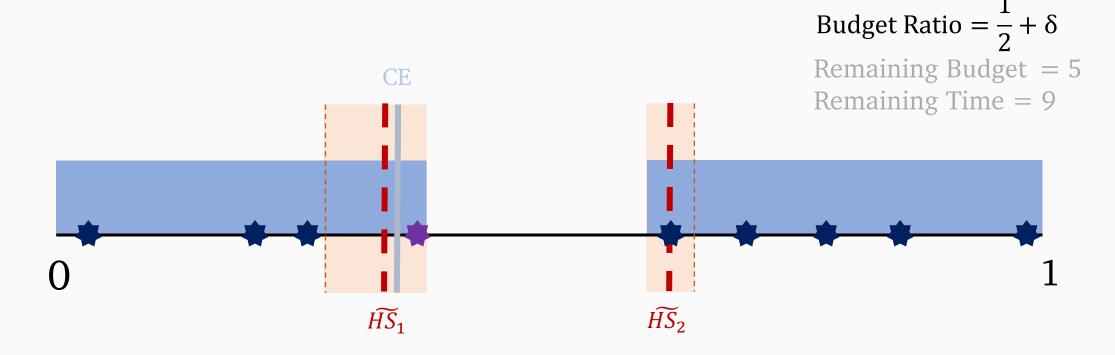


CwG and HS disagree, but expected cost of mistake $\sim \frac{1}{\text{remaining time}}$

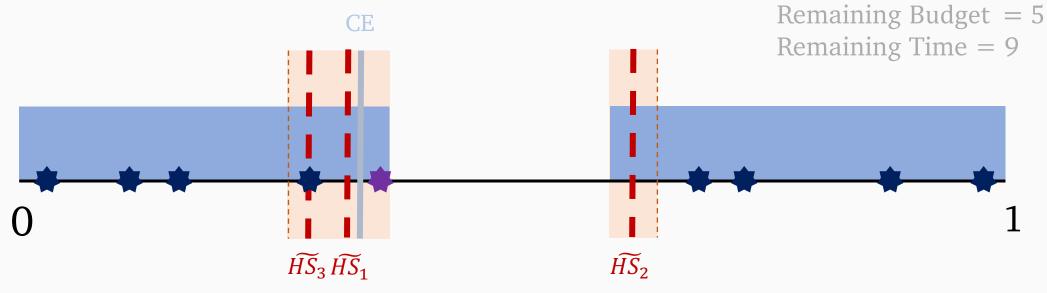
CwG via multiple simulations



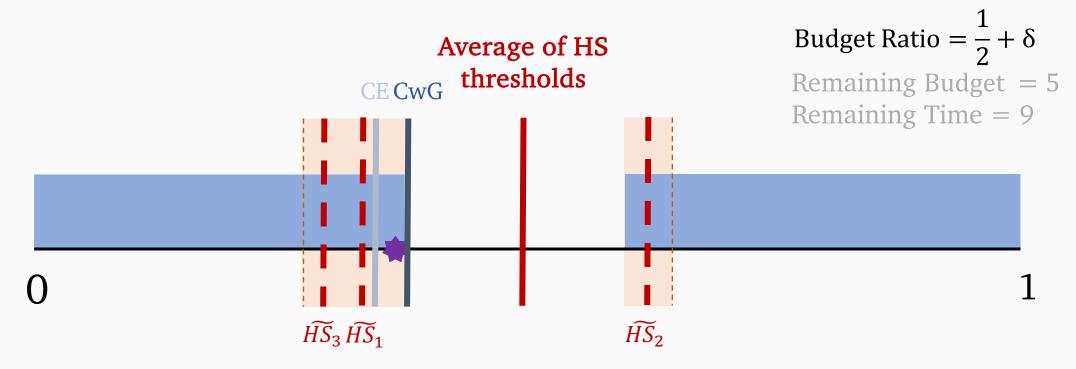
CwG via multiple simulations



CwG via multiple simulations Budget Ratio = $\frac{1}{2} + \delta$

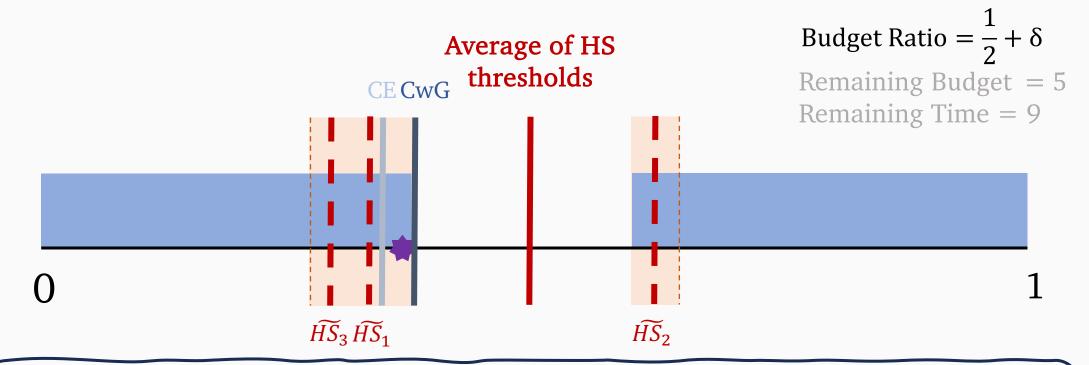


CwG via multiple simulations



same action under the CwG policy and using the average of the simulated HS thresholds

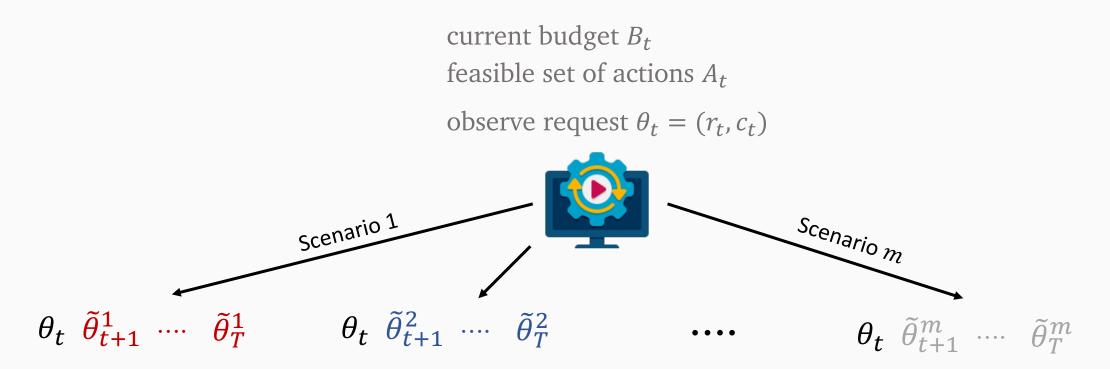
CwG via multiple simulations

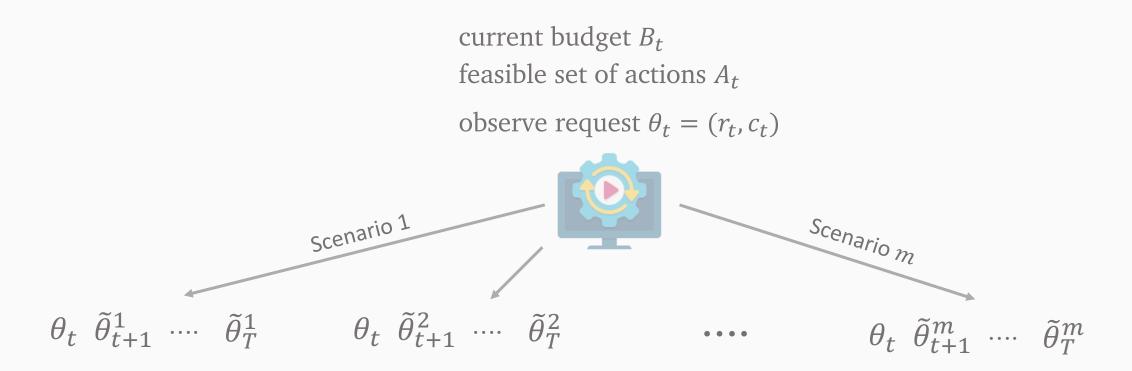


<u>Connection to "dual averaging" policy of Talluri and van Ryzin (1998)</u> think of the different HS thresholds as the shadow prices of the budget for different scenarios, the bid price is computed by averaging the HS thresholds

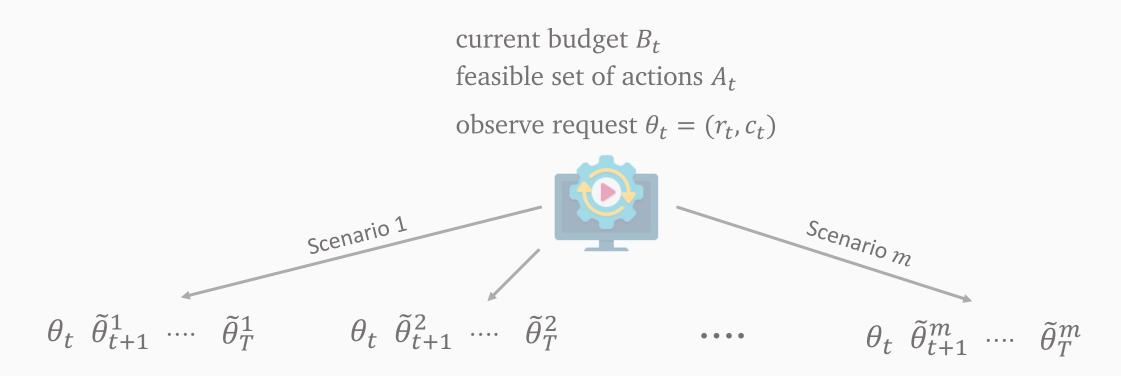
current budget B_t feasible set of actions A_t

observe request $\theta_t = (r_t, c_t)$



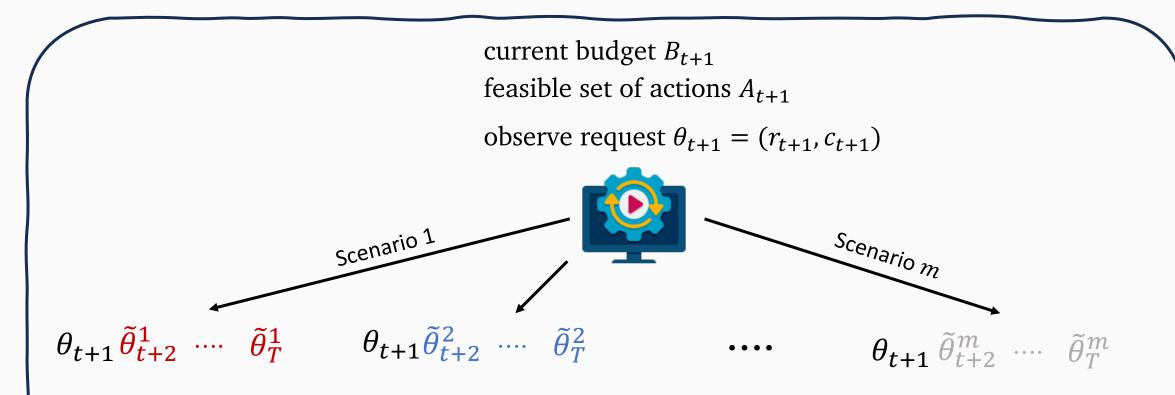


For each simulated scenario, compute the maximum total reward for each feasible current action at time t



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Take the action which maximizes the average total reward over the multiple simulated scenarios



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RAMS minimizes hindsightbased regret

Informal Meta Theorem.

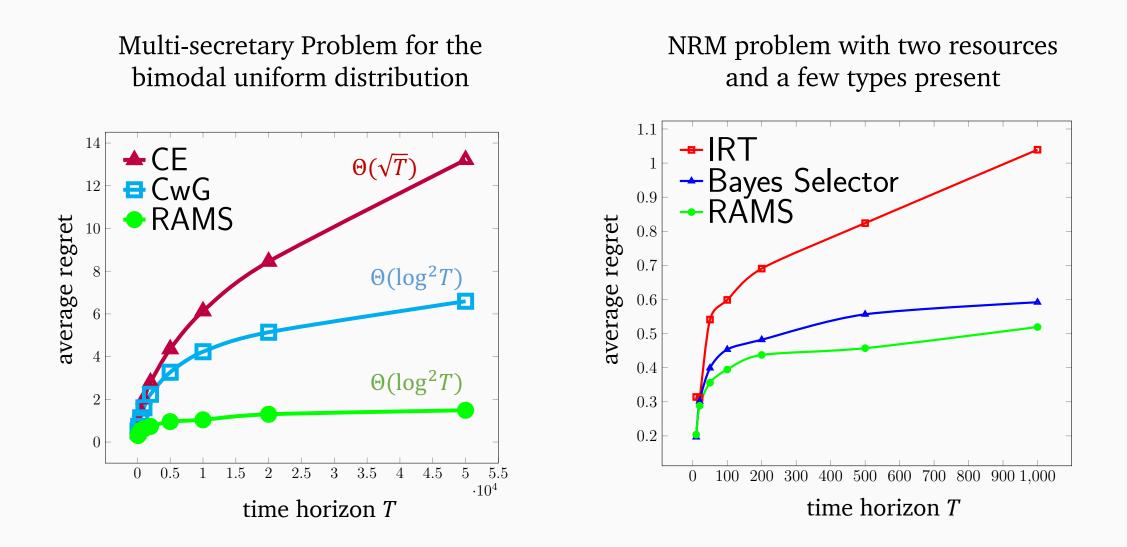
Given an dynamic resource allocation setting, if there exists an algorithm **ALG** satisfying certain technical conditions, then

Regret(RAMS) ≤ Regret Upper Bound of ALG + Sampling Error

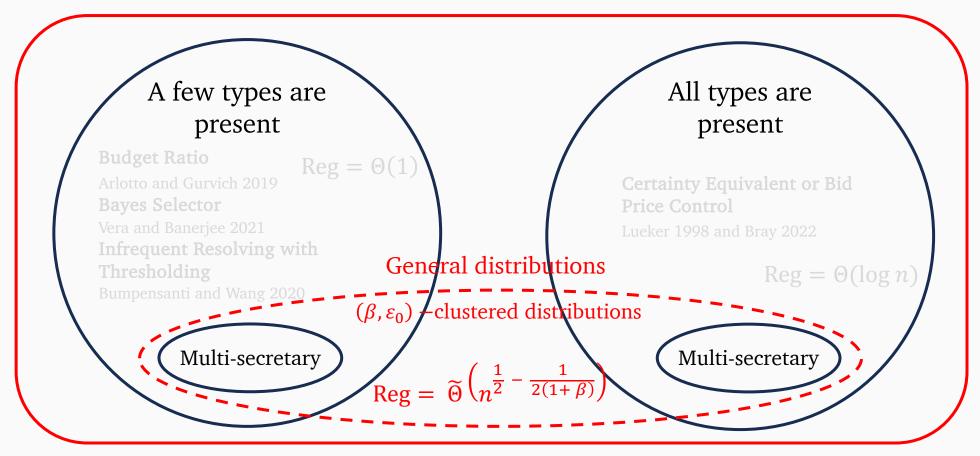
Corollary of the Meta Theorem

- **bounded regret** for NRM and online matching for a **few types** (Vera and Banerjee '21)
- logarithmic regret for NRM for many types with non-degeneracy assumption (Bray '22)
- <u>log-squared regret</u> for NRM for many types <u>without non-degeneracy</u> assumption (Jiang et. al '22)

Numerical Performance of RAMS



One policy to solve them all and with a simulator we bind them



Repeatedly Act using Multiple Simulations

Regret is with respect to the hindsight optimal problem

arxiv.org/abs/2205.09078

Thanks!